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 Subject: UAF in rsync 3.4.1 and below

## VULNERABILITY REPORT

=====

rsync: Use-After-Free via qsort Out-of-Bounds in receive\_xattr()  
 -----

Reporter: Przemyslaw Frasnunek <przemyslaw@...sunek.com>  
 Date: 2026-04-16  
 Affected: rsync 3.0.1 through 3.4.1 (all versions with xattr abbreviation)  
           Current development head (3.4.2dev, commit b905ab23) also affected.  
 File: xattrs.c, function receive\_xattr(), line 864  
 Severity: Medium / High  
 Attack vector: Network (malicious rsync sender)

### 1. SUMMARY

The receive\_xattr() function in rsync uses a wire-supplied count value as the length argument to qsort() instead of the actual number of items stored in the array. When xattr name filtering causes some items to be discarded, the count of items passed to qsort exceeds the number of valid items. The excess array positions contain stale data from a previously processed file's xattr list. After qsort reorders these stale entries into the stored positions, they are persisted in the global xattr list (rsync\_xal\_l) with dangling pointers to heap allocations that are subsequently freed and reallocated.

This creates use-after-free conditions in the receiver process: stale datum pointers are dereferenced in rsync\_xal\_set() for checksum comparison and in rcv\_xattr\_request() for name copying, and freed again in both functions, leading to double-free or free-of-allocated-memory scenarios.

### 2. AFFECTED CONFIGURATIONS

The bug is triggered when qsort runs over the stale array elements. The conditions depend on the operating system:

#### Linux:

The need\_sort flag defaults to 0 and is only set to 1 when a non-user namespace xattr (e.g. system.\*, security.\*) is received and prefixed with the rsync namespace prefix (RSYNC\_PREFIX = "user.rsync."). This prefixing only occurs when am\_root <= 0, specifically:

- rsync receiver running with --fake-super (am\_root = -1): VULNERABLE. Non-user namespace xattrs from the sender are prefixed and need\_sort is set. This is a common configuration for backup mirrors and rsyncd deployments that operate without root privileges.
- rsync receiver running as non-root (am\_root = 0) with xattr filter (--filter='x ...') configured: VULNERABLE. Non-user xattrs that pass the filter are prefixed and need\_sort is set.
- rsync receiver running as root without --fake-super: NOT VULNERABLE on Linux. Non-user xattrs are accepted without prefixing or sorting.
- rsync receiver running as non-root without xattr filter: NOT VULNERABLE. Non-user xattrs are silently discarded without setting need\_sort.

#### FreeBSD, macOS, and other non-Linux platforms:

The need\_sort flag defaults to 1 unconditionally (xattrs.c line 771).

The bug is triggerable whenever count > 1. All configurations running with -X/--xattrs are VULNERABLE.

### 3. ROOT CAUSE

In receive\_xattr() (xattrs.c), the wire-supplied xattr count for each file is read at line 786:

```
if ((count = read_varint(f)) != 0) {
    (void)EXPAND_ITEM_LIST(&temp_xattr, rsync_xa, count);
    temp_xattr.count = 0;
}
```

The EXPAND\_ITEM\_LIST call ensures the temp\_xattr.items array has room for count items. However, the loop that follows (lines 791-861) only adds items that pass namespace filtering and xattr filter checks. Each accepted item increments temp\_xattr.count via EXPAND\_ITEM\_LIST(..., 1). Items that are filtered out are freed and skipped via continue.

When some items are filtered, temp\_xattr.count < count. The array positions from temp\_xattr.count to count-1 are not cleared. Because temp\_xattr is a static variable reused across files, these positions contain stale rsync\_xa structs from the most recent file whose receive\_xattr() populated those indices.

At line 863-864:

```
if (need_sort && count > 1)
    qsort(temp_xattr.items, count, sizeof (rsync_xa), rsync_xal_compare_names);
```

qsort operates on count items, including the stale entries. Since qsort sorts by xattr name (rsync\_xal\_compare\_names), and the stale entries contain valid name pointers at this stage, the stale entries can be sorted into the first temp\_xattr.count positions of the array.

At line 866, rsync\_xal\_store(&temp\_xattr) copies exactly temp\_xattr.count items into the global rsync\_xal\_l list. If stale entries were sorted into positions 0..temp\_xattr.count-1, they are now permanently stored as the current file's xattr data, with datum and name pointers pointing into another file's xattr allocations.

### 4. EXPLOITATION

The stale rsync\_xa structs stored in rsync\_xal\_l reference datum buffers allocated during a prior file's receive\_xattr(). These buffers are freed when recv\_xattr\_request() processes that prior file's abbreviated xattr requests (line 753: free(old\_datum)), creating dangling pointers in the current file's stored xattr list.

The dangling pointers are subsequently dereferenced in two locations:

a) recv\_xattr\_request() for the current file (if the generator requests the stale items):

```
old_datum = rxa->datum; // dangling
rxs->datum = new_array(char, rxs->datum_len + rxs->name_len);
memcpy(name, rxs->name, rxs->name_len); // read-after-free
free(old_datum); // double-free / UAF free
```

b) rsync\_xal\_set() for the current file:

```
if (XATTR_ABBREV(rxsas[i])) { // struct field check
    ptr = get_xattr_data(fnamecmp, name, &len, 1); // read-after-free
    if (memcmp(sum, rxsas[i].datum + 1, ...) != 0) // read-after-free
        ...
    free(rxsas[i].datum); // double-free / UAF free
```

#### 4.1 Stock receiver exploitation constraints

Stock rsync does not write to freed memory after free(rxsas[i].datum) in rsync\_xal\_set(). Without a write-after-free primitive, redirecting tcache allocations to arbitrary addresses is not directly achievable.

The following primitives ARE available on stock rsync:

- Read-after-free: stale datum pointers are dereferenced for checksum comparison (memcmp against heap contents) and xattr name reads (memcpy from freed memory for name strings past tcache metadata offset).
- Free-of-allocated-memory: stale pointers that were re-allocated by another file's `recv_xattr_request` are freed again, bypassing glibc tcache key detection (the key field was overwritten by application data during re-allocation). This corrupts heap state by placing an in-use allocation on the tcache free list.
- Information disclosure: if `rsync_xal_set` processes a stale item via the non-abbreviated path (`lsetxattr`), tcache metadata (safe-linked pointers, tcache key) is written as xattr values on destination files. This leaks heap layout information to the filesystem.
- Denial of service: the double-free / heap corruption reliably crashes the receiver process.

## 5. TRIGGER CONDITIONS

For a malicious sender to trigger the vulnerability:

- a) The victim must run rsync with `-X (--xattrs)` to enable xattr transfer.
- b) On Linux, the victim must also use `--fake-super`, or have an xattr filter configured that passes non-user namespace xattrs. On FreeBSD/macOS, no additional flags are required.
- c) The sender must include at least one non-user namespace xattr in the triggering file's xattr list (Linux only, to set `need_sort=1`).
- d) The sender must set the wire count higher than the number of xattrs that will pass filtering, and arrange for a prior file in the transfer to populate the stale array positions with exploitable xattr structs.

The sender fully controls conditions (c) and (d) through the rsync protocol. Conditions (a) and (b) depend on the victim's rsync invocation.

## 6. SUGGESTED FIX

Replace `count` with `temp_xattr.count` in the `qsort` call at `xattrs.c` line 864:

```
Before:
    qsort(temp_xattr.items, count, sizeof (rsync_xa), rsync_xal_compare_names);
```

```
After:
    qsort(temp_xattr.items, temp_xattr.count, sizeof (rsync_xa),
rsync_xal_compare_names);
```

This ensures `qsort` only operates on actually populated items, preventing stale data from being sorted into the stored xattr list.

## 7. TIMELINE

- 2008-03-07 Bug introduced in commit `d724dd186` (rsync 3.0.1pre1).  
The commit added `qsort` to `receive_xattr()` for sorting xattrs after namespace prefix munging in `--fake-super` mode.
- 2026-04-16 This report.

## 8. REFERENCES

Source: <https://github.com/RsyncProject/rsync>  
Buggy commit: `d724dd186` ("Fixed the interaction of `--fake-super` with `--link-dest` & `--xattrs`. Fixed the munging of non-user namespace xattrs w/`--fake-super`. Fixed the sorting of received xattrs when name-munging occurs.")

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